

Lecture 0. A brief introduction of unbounded operators.

In this lecture, I'd like to introduce the basic notions about (unbounded) operators $A: D(A) \subseteq X \rightarrow Y$. Indeed, the most main result why we investigate the unbounded case for PDEs is that the differential operators are always unbounded.

Example 1. We consider the differentiation

$$A = \frac{d}{dt} : C^1 \subseteq C[a,b] \rightarrow C[a,b].$$

Then let $u_n(t) = \sin nt$, then $\|u_n\|_{C[a,b]} = 1$, but $\|A u_n\|_{C[a,b]} = n \rightarrow \infty$, which implies $A = \frac{d}{dt}$ is unbounded.

So we'd like to formulate a rigorous theory for the unbounded operators, which will serve as a fundamental language for PDE. It can be introduced briefly as following parts:

1) Basic notions: the extension and dual/adjoint of a unbounded operators.

The closed, symmetric and self-adjoint operators and its related facts.

Some applications to the Laplacian $-\Delta$, including its closure and self-adjoint extensions.

2) Spectral theory: The resolvent and spectrum for unbounded operators.

Two classification of spectrum set. spectrum of $-\Delta$.

Functional calculus: $A = \int_{\sigma(A)} \lambda dE(\lambda) \Rightarrow f(A) = \int_{\sigma(A)} f(\lambda) dE(\lambda)$.

Roughly saying, we can characterization the operators by its spectrum. And to obtain a kind of operator, we can operate directly on the spectrum and return the operation back to the operator level.

3) perturbation: We investigate the stability of the functional properties, including the closedness, invertibility, self-adjointness and the spectrum set

under a proper perturbation. Indeed, many kind of operators we concerned, such as Schrödinger operator $-\Delta + V$, Stokes operator $P\Delta$ and Lamé operator $-\Delta + \nabla(\nabla \cdot)$, can be viewed as perturbations of Laplacian $-\Delta$.

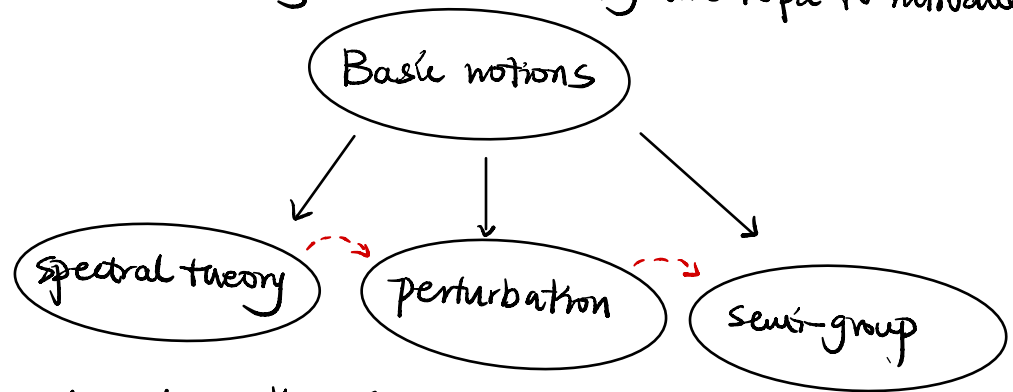
4) Semigroup theory: we rigorously define the semigroup e^{-tA} , and analyze when we can generate such kind of a semi-group.
 { Hille-Yosida, Zumer-Phillips, Stone

Using such concepts, we can strictly establish the well-posedness for abstract Cauchy problems: $\begin{cases} \partial_t u - Au = f \\ u(0) = u_0 \end{cases}$ (The Navier-Stokes and Schrödinger can be

reduced to such kind of abstract case.)

As for more deep topic I'd mentioned, it's better to introduce them after (pseudo-spectrum, analytic semigroups, spectral analysis) we finish the above fundament theorys.

I should mention that the basic notions in spectral / perturbation / semigroups are independent, are many books extract only one topic to introduce:



except some advanced results. I'd like to clarify this point for a clear sketch in mind.

As we introduce the differential operator $A = \frac{d}{dt}$, the domain of unbounded operators are not always everywhere-defined:

$$A: D(A) \subseteq X \rightarrow Y.$$

For example, the Laplacian $-\Delta: D \subseteq L^p \rightarrow L^p$ defined naturally in D .

But we may concern its extension in weak sense, (for example, in Sobolev spaces).

Definition 1. We define the graph of a operator A as:

$$P(A) = \{(x, Tx) \in X \times Y \mid x \in D(A)\}$$

Then we say \tilde{A} is a extension of A , denoted as $A \subseteq \tilde{A}$, if $P(A) \subseteq P(\tilde{A})$.

Remark. It is equivalent to say, for $x \in D(A)$, we have $x \in D(\tilde{A})$ and $Ax = \tilde{A}x$.

The another important concept is the dual (or conjugate, in some text book) of a unbounded operator (The analogue one is the adjoint for Hilbert case).

We recall that for bounded operator $A: X \rightarrow Y$, the dual operator $A': Y' \rightarrow X'$ is defined by following characterization:

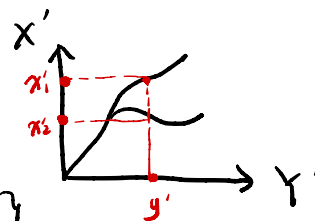
$$(Ax, y') = (x, A'y'), \quad \forall x \in X, y' \in Y'$$

The similar idea is applied to the unbounded operators, but a well-definedness argument should be noted due to not-everywhere-defined domain. And this is the reason why we always require dense-definedness when the dual arises.

Proposition 1. Let X, Y be N.V.S., and X', Y' be the dual.

Consider $A: D(A) \subseteq X \rightarrow Y$. Then if and only if A is densely defined, the following set can be viewed as a graph

$$P = \{(y', x') \in Y' \times X' \mid (y', Ax) = (x', x), \forall x \in D(A)\}$$



for some operators. In this case, we define $A': y' \mapsto x'$, $D(A') = P_{Y'}(P)$.

Remark. In the dense-defined case, A' exists and is characterized by

$$(y', Ax) = (A'y', x), \quad \forall x \in D(A), y' \in D(A'). \quad \star$$

Proof. That is to show, given if there exists $(y', x'_1), (y', x'_2)$ lays in P , then

$x'_1 = x'_2$. Indeed,

$$(x'_1, x) = (y', Ax) = (x'_2, x), \quad \forall x \in D(A), \Rightarrow (x'_1 - x'_2, D(A)) = 0.$$

Since $D(A)$ is dense in X , then $x'_1 = x'_2$.

Adjoint of a operator

Suppose $A : D(A) \subseteq H \rightarrow H$, where H is a Hilbert space. We can similarly define the adjoint of A if A is densely defined. And it is characterized by:

$$\langle Ax, y \rangle = \langle x, A^*y \rangle, \quad \forall x \in D(A), y \in D(A^*).$$

We show notice that the adjoint and the dual has the following relation:

$$A^* = b_H^{-1} \circ A' \circ b_H : D(A^*) = b_H^{-1} \circ D(A') \rightarrow H$$

Proposition. Suppose $A : D(A) \subseteq H \rightarrow H$, where H is a Hilbert space, then:

- 1) the dual A^* exists if and only if A is densely defined;
- 2) under the above case, if further $A^{**} = (A^*)^*$ exists, then $A \subseteq A^{**}$ naturally

Proof: 1) The idea is analogue to the dual case and so we omit it.

2) For $x \in D(A)$, we show that $x \in D(A^{**})$ as there exists unique $x^* = Ax$ ^{grant by existence} such that $\langle A^*y, x \rangle = \langle y, Ax \rangle = \langle y, x^* \rangle, \quad \forall y \in D(A^*).$

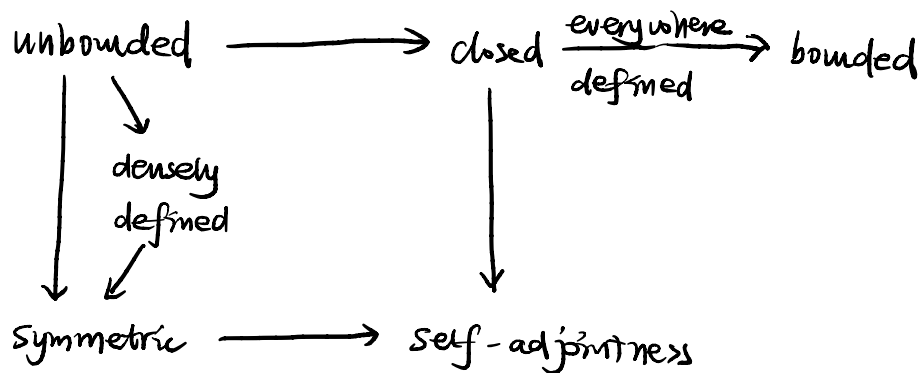
Consequently, we see $A^{**}x = (A^*)^*x = x^* = Ax$, which implies $A \subseteq A^{**}$.

Closedness, symmetry and the self-adjointness

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In this part, we'd like to introduce the notions about closedness and symmetry, and finally establish the basic facts about self-adjoint operators.

The mind trace can be formulated as following diagram:



First we'd like to define the relevant concepts about closedness:

Definition 2. Suppose $A: D(A) \subseteq X \rightarrow Y$. Then we say A is a closed operator if the graph $\overline{P(A)}$ is closed. More generally, A is closable if it has a closed extension. In this case, we define the smallest closed extension as the closure of A , denoted as \bar{A} .

Remark: There is a more useful definition to characterize the closability:

A is closable if and only if $\overline{P(A)}$ is a graph of some operator, in this case $\overline{P(A)} = P(\bar{A})$ exactly.

Indeed, we can show the equivalence:

\Leftarrow : if $\overline{P(A)} = P(S)$ for some operator S , then S must be a closed extension of A ;

\Rightarrow if S is a closed extension of A . Then by definition $\overline{P(A)} \subseteq P(S)$, which implies $\overline{P(A)}$ is uniquely correspondent ($\overline{P(A)}$) and thus a graph of some closed operator, denoted as R . The arbitrariness of S implies $R = \bar{A}$.

Some clarifications are listed below:

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Proposition. Suppose $A: D(A) \subseteq X \rightarrow Y$. Then

1) A is closed if and only if $\forall (x_n, Ax_n) \in P(A), (x_n, Ax_n) \rightarrow (x, y)$ implies $x \in D(A)$ and $y = Ax$.

2) A is closable if and only if $\forall (x_n, Ax_n) \in P(A), (x_n, Ax_n) \rightarrow (0, y)$ implies $y = 0$.

Proof: 1) Direct from the definition.

2) If A is closable, then the limit $(0, y)$ lays in $\overline{P(A)} = P(S)$ for some linear operator S , which implies $y = S(0) = 0$.

Conversely, we can define the operator S by following approximation:

$$Sx = \lim Ax_n, \quad \forall x \in \overline{D(A)},$$

there is well-defined by the assumption (free with choice of Ax_n).

Now we discuss the properties when it involves the adjoint.

Proposition: suppose $A: D(A) \subseteq H \rightarrow H$, where H is Hilbert:

1) iff A^* exists, then A^* is closed naturally;

2) iff further A^{**} exists, then A is closable. In this case, we have:

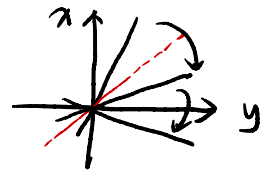
$$\bar{A} = A^{**}, \quad A^* = (\bar{A})^*$$

Indeed, we can observe from the characterization:

$$\langle Ax, y \rangle = \langle x, A^*y \rangle, \quad \forall x \in D(A),$$

$$\Leftrightarrow \langle (-Ax, x), (y, y^*) \rangle = 0, \quad \forall x \in D(A).$$

$$\Leftrightarrow (y, y^*) \in (J \circ P(A))^\perp, \quad J: (x, y) \mapsto (-y, x).$$



Then we can easily obtain the following lemma:

Lemma: T^* exists if and only if $(J \circ P(A))^\perp$ is a graph of some operator. In this case, we have $(J \circ P(A))^\perp = P(T^*)$.

Proof of proposition: 1) It is directly from the lemma, since complement is always closed. ⁷

2) we should notice that: $\overline{P(A)} = (P(A)^\perp)^\perp = (\underbrace{J \circ J \circ P(A)^\perp}_{\text{invariant for linear space}})^\perp = (J \circ P(A^*))^\perp$
 (if A^{**} exists) = $P(A^{**})$. invariant for linear space ↕

This implies, A is closable, and its closure is exactly A^{**} . ($\bar{A} = A^{**}$)

and $(\bar{A})^* = (A^{**})^* = (A^*)^{**} = \bar{A}^*$.

Converse side check later. Δ .

Some other properties will be checked later when we need.

Theorem. (Closed Graph). Suppose $A: D(A) \subseteq X \rightarrow Y$ is closed, then A is bounded iff e.d.

Theorem. (Hille). $A \int = \int A$.

Method to compute a closure.

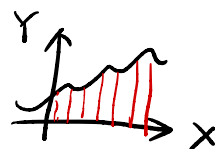
Here we'd like to give a method to compute the closure extension, and take the Laplacian $-\Delta: D \subseteq L^p \rightarrow L^p$ for example.

Proposition: Suppose X, Y is Banach, then $A: D(A) \subseteq X \rightarrow Y$ is closed if and only if $D(A)$ is complete under the graph norm $\|\cdot\|_P$ defined by

$$\|x\|_P = \|(x, Ax)\|_{X \times Y} = \|x\|_X + \|Ax\|_Y.$$

Proof: we consider the isometry:

$$\begin{array}{ccc} \Phi: D(A) & \rightarrow & P(A), \quad x \mapsto (x, Px) \\ \downarrow & & \downarrow \\ \|\cdot\|_P & & \|\cdot\|_{X \times Y} \end{array}$$



$P(A)$ is closed in $X \times Y \iff P(A)$ is complete in $X \times Y \iff D(A)$ is complete in $\|\cdot\|_P$.

Then the closure of $-\Delta: D \subseteq L^2 \rightarrow L^2$ has the domain:

$$D(\bar{-\Delta}) = \overline{D(-\Delta)P} = \overline{D(-\Delta)H^2} = \bar{D}H^2 = H_0^2.$$

Since $\|u\|_P = \|u\|_{L^2} + \|\Delta u\|_{L^2} \underset{u \in D}{\sim} \|u\|_{H^2}$ by Poincaré inequality.

Symmetry and self-adjointness.

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As the adjoint notions are always involved in this section, we'd like to assume $A: D(A) \subseteq H \rightarrow H$ is densely defined first.

Definition: We say $A: D(A) \subseteq H \rightarrow H$ is symmetric if $A \subseteq A^*$, self-adjoint if $A = A^*$. And say A is essentially self-adjoint if \bar{A} is self-adjoint.

We'd like to clarify these notions in later discussion. Indeed, a symmetric operator A can be characterized by:

$$\langle Ax, y \rangle = \langle x, Ay \rangle, \quad \forall x \in D(A), y \in D(A).$$

Then adjoint relation of a symmetric is quite interesting:

Proposition: Suppose A is a ^(closed) symmetric operator. Then A^* , A^{**} exists and $A \subseteq A^{**} \subseteq A^*$. ⁽⁼⁾

Proof: A^{**} exists since $A \subseteq A^*$ implies $D(A^*) \supseteq D(A)$ is dense. In this case, we know that \checkmark A is closable and $A \subseteq \bar{A} = A^{**}$. Moreover, since A^* is closed, then $\bar{A} \subseteq A^*$, which grant finally $A \subseteq \bar{A} \subseteq A^*$.

$$\begin{array}{c} A \subseteq \bar{A} \subseteq A^* \\ \parallel \\ A^{**} \end{array}$$

The following proposition give several equivalent condition when a symmetric operator is self-adjoint.

Proposition. Suppose $A: D(A) \subseteq H \rightarrow H$ is a symmetric operator, then the following statements are equivalent:

- 1) A is self-adjoint; 2) $D(A) = D(A^*)$; 3) A is closed and A^* is symmetric;
- 4) A is closed and $\text{Ker}(A \pm i) = \{0\}$; 5) $\text{Ran}(A \pm i) = X$.
- 6) $\sigma(A) \subseteq \mathbb{R}$. (prove later).

Proof: $1) \Leftrightarrow 2)$ is clearly from the definition.

$1) \Leftrightarrow 3)$ is from the fact: $A^* \subseteq (A^*)^* = \bar{A} = A$.

$1) \Rightarrow 4)$ Self-adjoint implies $A = A^*$ is closed clearly, moreover, $\forall x \in \ker(A^* \pm i)$

$$(A^* \pm i)x = 0,$$

$$F_i \langle x, x \rangle = \langle F_i x, x \rangle = \langle Ax, x \rangle = \langle x, Ax \rangle = \langle x, F_i x \rangle = \pm i \langle x, x \rangle$$

$$\Rightarrow \langle x, x \rangle = 0 \Rightarrow x = 0.$$

$4) \Rightarrow 5)$ clearly $\overline{\text{Ran}(T \pm i)} = X$, we show that $\text{Ran}(T \pm i)$ is closed.

The important observation is that

$$\begin{aligned} \|(T \pm i)x\|^2 &= \langle (T \pm i)x, (T \pm i)x \rangle \\ &= \langle Tx, Tx \rangle + \langle x, x \rangle + \underbrace{\langle Tx, \pm i x \rangle + \langle \pm i x, Tx \rangle}_{=0} \\ &= \|Tx\|^2 + \|x\|^2. \end{aligned}$$

Then for $\{(T \pm i)x_n\}$ Cauchy, so is $\{x_n\}$ and $\{Tx_n\}$ and closedness

ensure $x_n \rightarrow x_0$ and $Tx_n \rightarrow Tx_0$, which implies: $(T \pm i)x_n \rightarrow (T \pm i)x_0$.

$5) \Rightarrow 1)$. We need to show for $x \in D(A^*)$, $x \in D(A)$. Since $\text{Ran}(A \pm i) = X$, we can find $x' \in D(A)$ such that $(A \pm i)x' = (A^* \pm i)x$. Following we show $x = x'$:

$$\triangle A \subseteq A^* \Rightarrow (A \pm i)x' = (A^* \pm i)x \Rightarrow (A^* \pm i)(x' - x) = 0 \Rightarrow x' - x = 0.$$

A direct corollary from 4), 5) is the following characterization for essentially self-adjoint:

Corollary: A symmetric operator $A: D(A) \subseteq H \rightarrow H$ is essentially self-adjoint iff

$$1) \ker(A^* \pm i) = 0; \text{ or } 2) \overline{\text{Ran}(A \pm i)} = H.$$

